

Complex networks

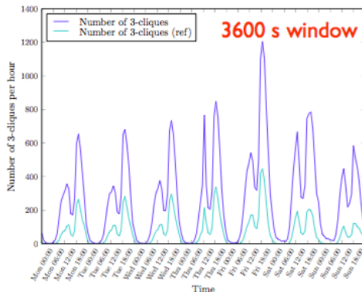
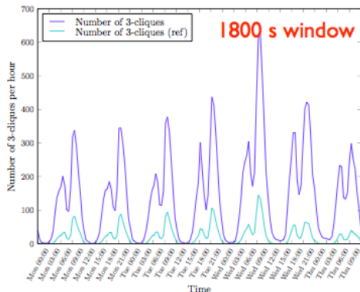
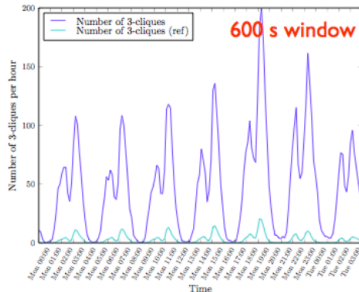
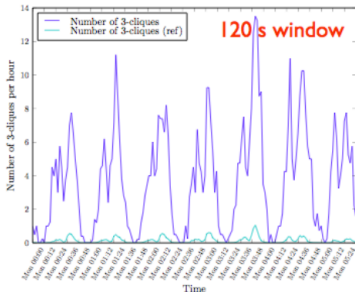
Diffusion and spreading on networks

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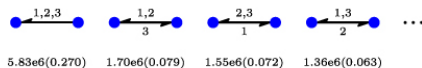
April 26, 2023

Temporal motifs: occurrence of triangles

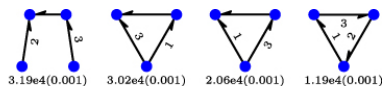


Temporal motifs: occurrence of ordered sequences

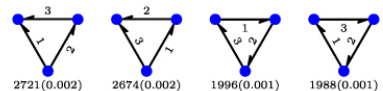
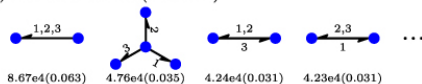
Most frequent ones



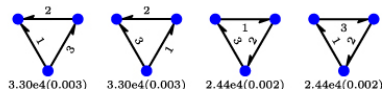
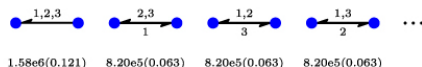
Least frequent ones



(b) TIME-SHUFFLED (unbiased)



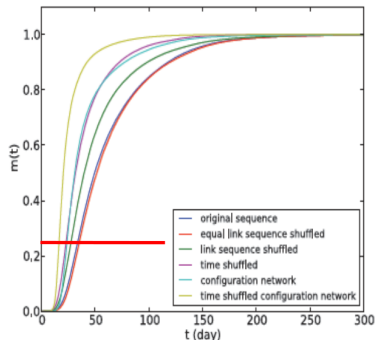
(c) TIME-SHUFFLED ($m = 32$)



Importance of different effects in temporal spreading

- ▶ Equal-weight link-sequence shuffled: Whole single-link event sequences are randomly exchanged between links having the same number of events
- ▶ Only link-link correlation is destroyed

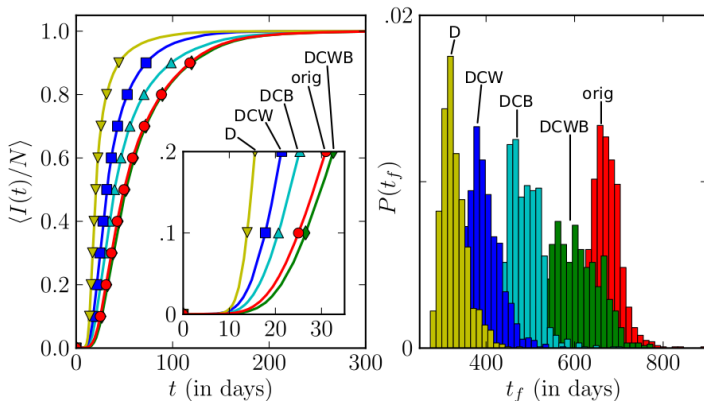
Event sequence	D	C	W	B	E	25%
Original	✓	✓	✓	✓	✓	33.7
Config. shuffle	✓	×	×	×	×	16.4
Config. keep	✓	×	×	✓	×	23.8
Orig. shuffle	✓	✓	✓	×	×	22.9
Shuffle. keep	✓	✓	×	✓	×	27.5
W keep sh.,keep	✓	✓	✓	✓	×	35.3



Importance of different effects in temporal spreading

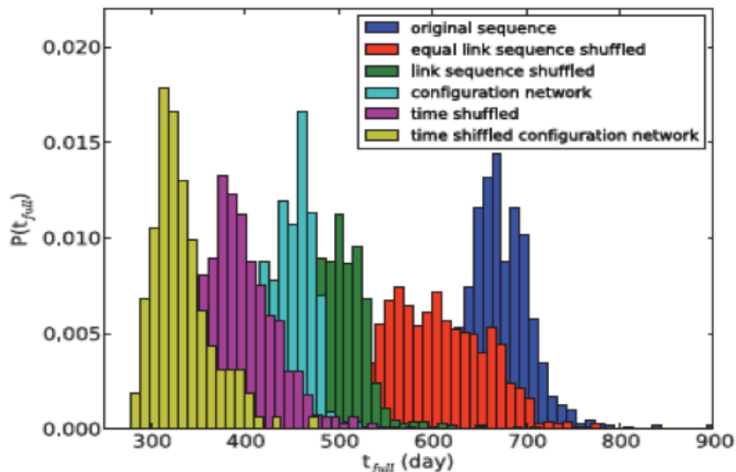
► Long time behaviour:

Event sequence	D	C	W	B	E	25%
Original	✓	✓	✓	✓	✓	33.7
Config. shuffle	✓	×	×	×	×	16.4
Config. keep	✓	×	×	✓	×	23.8
Orig. shuffle	✓	✓	✓	×	×	22.9
Shuffle. keep	✓	✓	×	✓	×	27.5
W keep sh.,keep	✓	✓	✓	✓	×	35.3



Importance of different effects in temporal spreading

- ▶ Everything slows down the spreading
- ▶ Burstiness has higher impact than topological structures

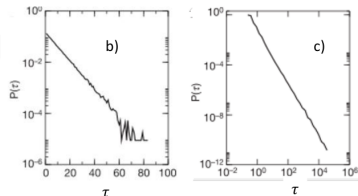


Interevent time

- ▶ Time interval between successive events τ
- ▶ Distribution of τ is $P(\tau)$
- ▶ Distribution is characterized by the average $\langle\tau\rangle$ and the variance σ
- ▶ Burstiness:

$$B = \frac{\sigma - \langle\tau\rangle}{\sigma + \langle\tau\rangle}$$

- ▶ (a) $B = -1$: deterministic, (b) $B = 0$: Poisson, (c) $B = 1$: bursty



Spreading on networks

- ▶ One of the most important problems on networks
- ▶ Also one of the real success
- ▶ This lecture:
 - ▶ Advanced mean-field calculations
 - ▶ Cascade models
 - ▶ Spreading in temporal networks

Epidemic models: notations

► States:

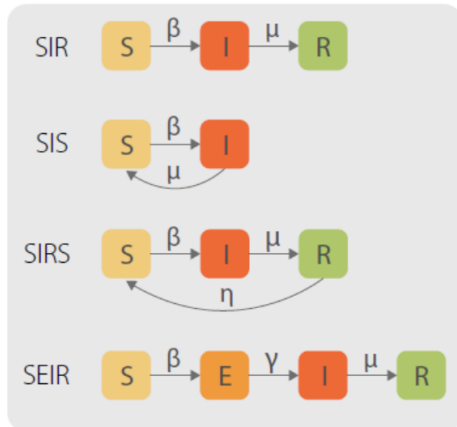
► S: susceptible

► I: Infected

► R: Recovered (immune)

► E: Exposed (infected but not yet infecting)

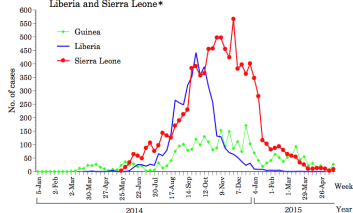
► Rates: β , μ , η , γ



SIR reality vs. model

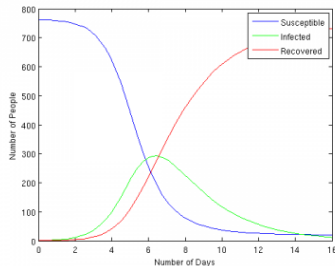
- ▶ Perfect mixing
- ▶ Everybody can meet everybody
- ▶ Ebola

Figure 1. Confirmed weekly Ebola virus disease cases reported from Guinea, Liberia and Sierra Leone*



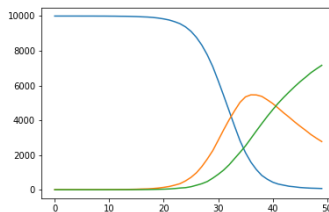
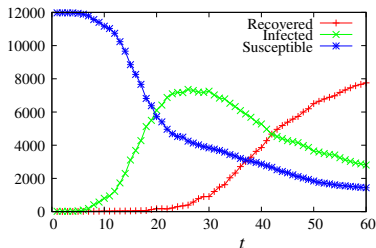
* WHO Ebola Situation Report - 20 May 2015 (<http://apps.who.int/ebola/en/ebola-situation-reports>)
Guinea (Patient database), Liberia (Situation report, as of 6 May 2015), Sierra Leone (Patient database up to 1 May 2015, Situation report from 8 March 2015)

IASR



SIR reality vs. model

- ▶ Perfect mixing
- ▶ Everybody can meet everybody
- ▶ Covid-19, South Korea
- ▶ Susceptible approximated



SIS model: mean field

- ▶ Perfect mixing
- ▶ Everybody can meet everybody
- ▶ The different type meet with probability proportional to their density
- ▶ Density of types:

$$\rho^\alpha = N^\alpha / N$$

- ▶ The mean field SIR equations:

$$\begin{aligned}\frac{d\rho^I}{dt} &= \beta\rho^I\rho^S - \mu\rho^I \\ \frac{d\rho^S}{dt} &= -\beta\rho^I\rho^S + \chi\rho^I\end{aligned}$$

where $\chi = \mu$ for SIS and $\chi = 0$ for SIR.

Epidemic threshold

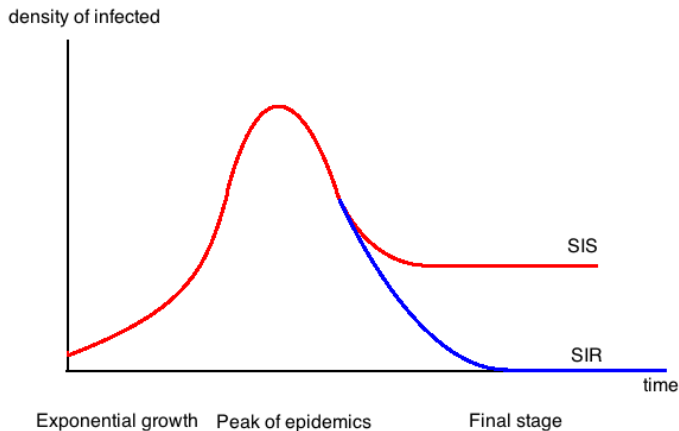
- ▶ Linearization: $\rho^I \ll 1$, $\rho^S \simeq 1$

$$\begin{aligned}\frac{d\rho^I}{dt} &= \beta\rho^I\rho^S - \mu\rho^I \\ \frac{d\rho^I}{dt} &\simeq (\beta - \mu)\rho^I \\ \rho^I(t) &\simeq \rho^I(0) \exp[(\beta - \mu)t]\end{aligned}$$

- ▶ Two regimes:
 - ▶ $\beta < \mu$: Disease dies out
 - ▶ $\beta > \mu$: Disease spreads
- ▶ Reproduction number: $R_0 = \beta/\mu$
- ▶ The epidemic threshold for perfectly mixing population is $R_0 = 1$ above which the epidemic spreads.

SIR/SIS

- ▶ Above the epidemic threshold
- ▶ In SIS dynamic equilibrium



SIS: Individual based mean field (IBMF)

- ▶ Markov chain approach
- ▶ Two state $X_i = 1$ for I and $X_i = 0$ for S .
- ▶ $E[X_i(t)]$ expected value of X_i
- ▶ a_{ij} element of the adjacency matrix
- ▶ The Master equation:

$$\frac{dE[X_i(t)]}{dt} = E \left[-\mu X_i(t) + (1 - X_i(t))\beta \sum_j a_{ij} X_j(t) \right]$$

- ▶ Introducing $\lambda = \beta/\mu$ and rescaling the time by $1/\mu$
- ▶ For static network:

$$\frac{d\rho_i^I(t)}{dt} = -\rho_i^I(t) + \lambda \sum_j a_{ij} \rho_j^I(t) - \lambda \sum_j a_{ij} E[X_i(t)X_j(t)]$$

SIS: Individual based mean field (IBMF)

- ▶ Markov chain approach
- ▶ The Master equation for static network:

$$\frac{d\rho_i^I(t)}{dt} = -\rho_i^I(t) + \lambda \sum_j a_{ij} \rho_j^I(t) - \lambda \sum_j a_{ij} E[X_i(t)X_j(t)]$$

- ▶ No explicit solution due to the two term correlations $E[X_i(t)X_j(t)]$.
- ▶ Joint probability distribution cannot be calculated
- ▶ Assumption: neighboring nodes are statistically independent:

$$E[X_i(t)X_j(t)] \equiv E[X_i(t)]E[X_j(t)] = \rho_i^I(t)\rho_j^I(t)$$

- ▶ The Master equation for the SIS model thus reads:

$$\frac{d\rho_i^I(t)}{dt} = -\rho_i^I(t) + \lambda[1 - \rho_i^I(t)] \sum_j a_{ij} \rho_j^I(t)$$

SIS: Individual based mean field (IBMF)

- The Master equation for static network SIS model in the independent neighbors limit:

$$\frac{d\rho_i^I(t)}{dt} = -\rho_i^I(t) + \lambda[1 - \rho_i^I(t)] \sum_j a_{ij} \rho_j^I(t)$$

- Loss term: probability that node i is infected times the rate of recovery (hidden in the rescaled time)
- Gain term: probability that node i is **susceptible**, times **the total probability that any of its nearest neighbors is infected**, times **the effective transmission rate $\lambda = \beta/\mu$**

SIS: Individual based mean field (IBMF)

- ▶ The Master equation for static network SIS model in the independent neighbors limit:

$$\frac{d\rho_i^I(t)}{dt} = -\rho_i^I(t) + \lambda[1 - \rho_i^I(t)] \sum_j a_{ij} \rho_j^I(t)$$

- ▶ Linear stability analysis

$$\frac{d\rho_i^I(t)}{dt} \simeq -\rho_i^I(t) + \lambda \sum_j a_{ij} \rho_j^I(t) = \sum_j J_{ij} \rho_j^I(t)$$

with $J_{ij} = -\delta_{ij} + \lambda a_{ij}$

- ▶ An endemic state occurs when Λ_1 the largest eigenvalue of J is positive.
- ▶ The epidemic threshold is thus:

$$\lambda > \lambda_c^{\text{IBMF}} \equiv \frac{1}{\Lambda_1}$$

SIS: Individual based mean field (IBMF)

- ▶ For networks with power law degree distribution $P(k) \sim k^{-\gamma}$
- ▶ Largest eigenvalue:

$$\Lambda_1 = \min(\sqrt{k_{\max}}, \langle k^2 \rangle / \langle k \rangle)$$

- ▶ The epidemic threshold:

$$\lambda_c^{\text{IBMF}} = \begin{cases} \frac{1}{\sqrt{k_{\max}}} & \text{if } \gamma \geq 5/2 \\ \frac{\langle k \rangle}{\langle k^2 \rangle} & \text{if } 2 < \gamma < 5/2 \end{cases}$$

- ▶ In both cases $\lim_{N \rightarrow \infty} \lambda_c^{\text{IBMF}} = 0$ for scale free networks
- ▶ Of course in finite system there is a small deviation but in infinite systems there is no epidemic threshold since the largest degree is infinite.

SIS: Degree based mean field (DBMF)

- ▶ All nodes with the same degree are statistically equivalent.
- ▶ Degree has a maximum k_{max}
- ▶ Number of equations k_{max}
- ▶ Conditional probabilities: $P(k'|k)$ probability that a node with degree k is connected to a node of degree k'
- ▶ $P(k'|k)$ is the same for all k degree nodes.
- ▶ In the case of uncorrelated networks:

$$P(k'|k) = \frac{k'P(k')}{\langle k \rangle}$$

SIS: Degree based mean field (DBMF)

- ▶ $\rho_k^I(t)$ is the probability that a node of degree k is infected at time t
- ▶ Master equation:

$$\frac{d\rho_k^I(t)}{dt} = -\rho_k^I(t) + \lambda k [1 - \rho_k^I(t)] \sum_{k'} P(k'|k) \rho_{k'}^I(t)$$

- ▶ Note that the factor k in the gain term is for the number of links the node of degree k has with that chance to get infected
- ▶ Linearized version

$$\frac{d\rho_k^I(t)}{dt} \simeq -\rho_k^I(t) + \lambda k \sum_{k'} P(k'|k) \rho_{k'}^I(t) = \sum_{k'} J_{kk'} \rho_{k'}^I(t)$$

- ▶ With

$$J_{kk'} = -\delta_{kk'} + \lambda k P(k'|k)$$

SIS: Degree based mean field (DBMF)

- ▶ Linearized version

$$\frac{d\rho_k^I(t)}{dt} \simeq -\rho_k^I(t) + \lambda k \sum_{k'} P(k'|k) \rho_{k'}^I(t) = \sum_{k'} J_{kk'} \rho_{k'}^I(t)$$

- ▶ With

$$J_{kk'} = -\delta_{kk'} + \lambda k P(k'|k)$$

- ▶ There is an epidemic state if

$$\lambda > \lambda_c^{DBMF} = \frac{1}{\Lambda_1}$$

where again Λ_1 is the largest eigenvalue of J

SIS: Degree based mean field (DBMF)

- ▶ The epidemic threshold, for uncorrelated networks with

$$P(k'|k) = \frac{k'P(k')}{\langle k \rangle}$$

- ▶ Probability to find an infected node following a randomly chosen edge

$$\Theta = \sum_{k'} P(k'|k) = \frac{k'P(k')}{\langle k \rangle} \rho_{k'}^I(t)$$

The Master equation of the Degree based mean field is

$$\frac{d\rho_k^I(t)}{dt} = -\rho_k^I(t) + \lambda k[1 - \rho_k^I(t)]\Theta$$

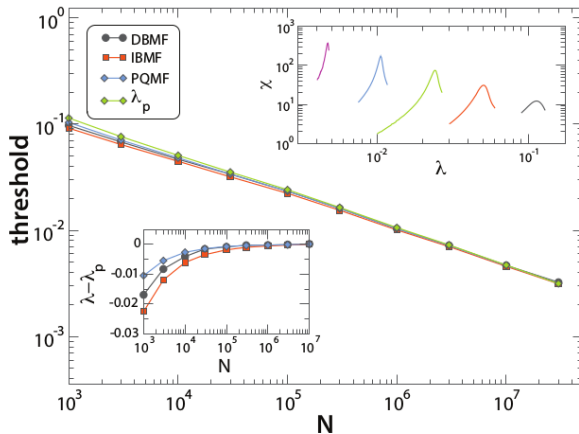
- ▶ The two latter equations can be solved in self-consistently.

SIS: Degree based mean field (DBMF)

- ▶ The self-consistent solution allows for an epidemic state only if

$$\lambda > \lambda_c^{DBMF} = \frac{\langle k \rangle}{\langle k^2 \rangle}$$

- ▶ For power law degree distribution with exponent $2 < \gamma \leq 3$
The threshold is 0 in the infinite limit.



SIS: Comparison

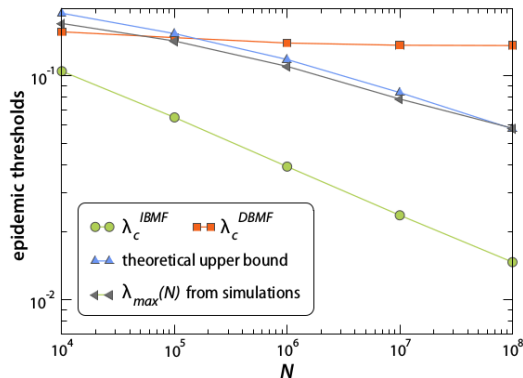
► IBMF

$$\lambda_c^{\text{IBMF}} = \begin{cases} \frac{1}{\sqrt{k_{\max}}} & \text{if } \gamma \geq 5/2 \\ \frac{\langle k \rangle}{\langle k^2 \rangle} & \text{if } 2 < \gamma < 5/2 \end{cases}$$

► DBMF

$$\lambda > \lambda_c^{\text{DBMF}} = \frac{\langle k \rangle}{\langle k^2 \rangle}$$

► Network: scale free with $\gamma = 2.25$

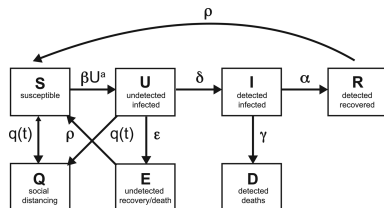


SIS: Epidemic threshold

- ▶ It seems that in the scale free networks in the infinite system there is no epidemic threshold
- ▶ Numerical simulations show also this picture
- ▶ Note that in the SIS model there is a dynamic steady state with a fraction of infected nodes
- ▶ In scale-free networks only part of the system will be infected, the hubs and the immediate neighborhood.
- ▶ Concepts of:
 - ▶ **Epidemic state:** Homogeneously infected
 - ▶ **Active state:** Small finite active part
- ▶ Thankfully real systems are never infinite and never scale-free

Problems with SIR model

- ▶ Fully mixed society \rightarrow network
- ▶ Disease either dies out fast or infects the whole society even on networks \rightarrow geographical location with travel links effected by the infection
- ▶ Inhomogeneous society \rightarrow age groups and connections as in SBM
- ▶ Computational limitations \rightarrow two level systems
- ▶ Government measures \rightarrow new states in the SIR model



Khan, Van Bussel, Hussain, Epidemiology and Infection 2020

Immunization

- ▶ Often the task is to stop the spreading
- ▶ Sometimes one can immunize part of the society
- ▶ Can we stop the spreading?
- ▶ Example:
 - ▶ Of course, if every newborn baby is vaccinated, the population is safe. This is the way, how smallpox (Variola) was defeated.
 - ▶ Estimated death in 20 th century: 300 Million
 - ▶ Estimated infected in 1967: 15 Million
 - ▶ 1979: WHO declared smallpox eradicated

Immunization

- ▶ Epidemic threshold (complete graph/fully mixed state):

$$R_0 = \frac{\beta}{\mu} \begin{cases} > 1 & \text{outbreak} \\ = 1 & \text{threshold} \\ < 1 & \text{localized} \end{cases}$$

- ▶ The density of the immune vertices is g , then:

$$\beta' = \beta(1 - g)$$

- ▶ The threshold for networks

$$\frac{\beta(1 - g)}{\mu} = \frac{\langle k \rangle}{\langle k^2 \rangle}$$

- ▶ For infinitely large scale free network with $\gamma \leq 3$ we get $g_c = 1$
- ▶ For random immunization everybody must be vaccinated

Immunization

- ▶ Epidemic threshold for networks

$$\frac{\beta(1-g)}{\mu} = \frac{\langle k \rangle}{\langle k^2 \rangle}$$

- ▶ Targeted immunization: immunize high degree nodes
- ▶ This decreases the variance faster than the average

$$\frac{\langle k \rangle_g}{\langle k^2 \rangle_g} > \frac{\beta(1-g)}{\mu}$$

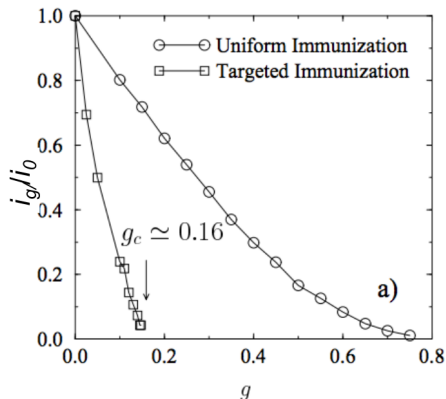
which defines the critical value of g

Immunization

- Targeted immunization: immunize high degree nodes

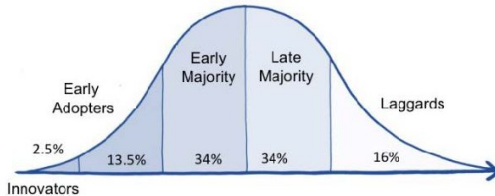
$$\frac{\langle k \rangle_g}{\langle k^2 \rangle_g} > \frac{\beta(1-g)}{\mu}$$

which defines the critical value of g

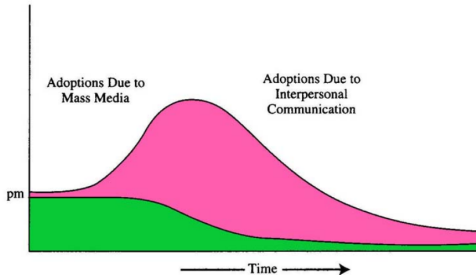


Innovation spreading

► Rogers (1962)



► Mahajan, Muller and Bass (1990)



Threshold model

- ▶ Sometimes the spreading is due to load from the neighbors
- ▶ E.g. if too many of my neighbors are infected I will also get infected
- ▶ Innovation spreading: many of my friends have iPhone I will also get one.



Threshold model

- ▶ Networks with average degree $\langle k \rangle = z$
- ▶ Nodes have threshold ϕ_i
- ▶ If the number of active nodes in the neighborhood reach ϕ_i then the node becomes active (too many friends have some product I will also buy it)
- ▶ Start from a small seed
- ▶ If thresholds are sufficiently low cascades may propagate through the whole system (size $\sim \mathcal{O}(N)$)

Watts, A simple model of global cascades on random networks (2002)

Threshold model

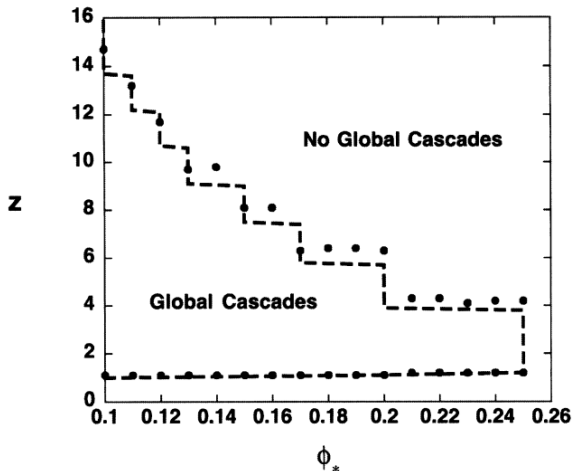
- ▶ In large uncorrelated random networks there are hardly any triangles
- ▶ *Vulnerable* nodes are the ones where the threshold is less than $\phi_i < 1/k_i$, one neighbor is enough to get infected
- ▶ Global cascade is possible if these nodes percolate
- ▶ This is the *cascade condition*

$$z > \sum_k k(k-1)P(k)P(\phi \leq 1/k)$$

- ▶ $k(k-1)$ increases with k
- ▶ $P(\phi \leq 1/k)$ decreases with k
- ▶ Two or 0 solutions

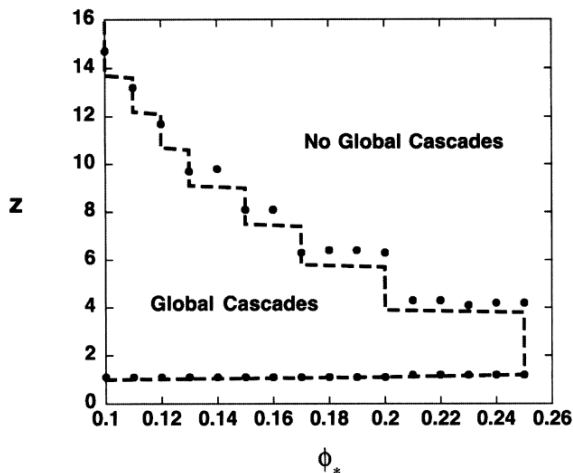
Threshold model: Phase diagram

- Points simulation
- Dashed line calculated threshold



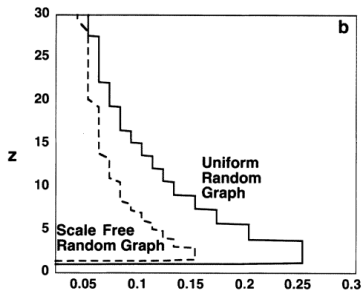
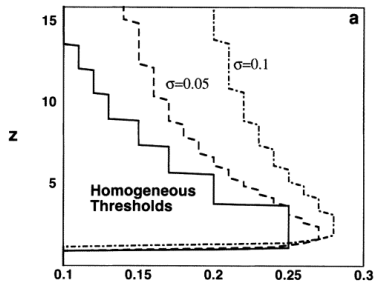
Threshold model: Phase diagram

- ▶ Top line: first order phase transition of cascades
- ▶ Bottom line: second order phase transition of network percolation limit



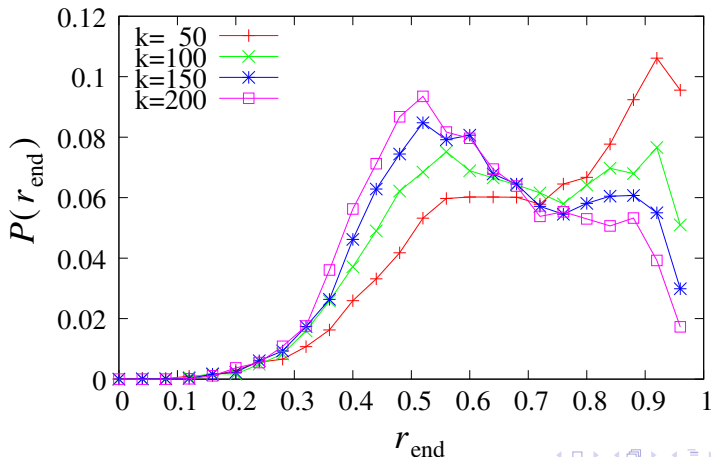
Threshold model: Phase diagram

- ▶ ϕ With normal distribution and σ variance
- ▶ Scale free graph



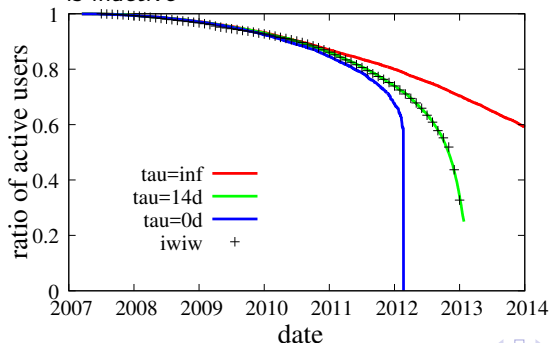
Fall of a social network site

- ▶ Users leave due to exogenous effects (advertisements, news, etc.)
- ▶ Users leave if some part of their friends leave.
- ▶ This depends on the embeddedness of the user



Fall of a social network site: Model

- ▶ Users leave due to exogenous effects (advertisements, news, etc.):
 - ▶ Here rate of leave increases with time as was the popularity of the alternative site
 - ▶ Users with low degree are more susceptible to global effects
- ▶ Users leave if their friends leave.
 - ▶ Threshold model with threshold above 50%
 - ▶ Leave is not immediate one needs time (τ) to recognize friend is inactive



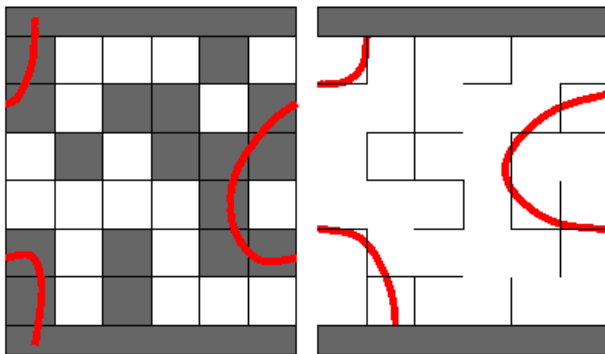
Percolation



Percolation

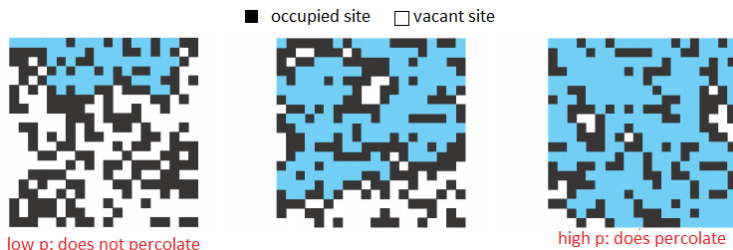
Behavior of **connected** cluster

- ▶ Site percolation
- ▶ Bond percolation



Percolation model

- ▶ Random environment
- ▶ With probability p site vacant (conducts)
- ▶ Two states: percolates or not!
- ▶ Percolation: presence of infinite cluster, in infinitely large system the cluster holds finite fraction of nodes.



Percolation theory

Questions (in infinite systems):

1. Is there an infinite cluster in infinite systems?
2. How many infinite clusters are there?
3. Mean cluster size (without the infinite one)?
4. Cluster size distribution

Answers:

1. Above a critical density with probability 1 below it with probability 0
2. Only 1!
3. Decreases as a power law away from the critical density
4. Power law

Percolation theory

Questions (in infinite systems):

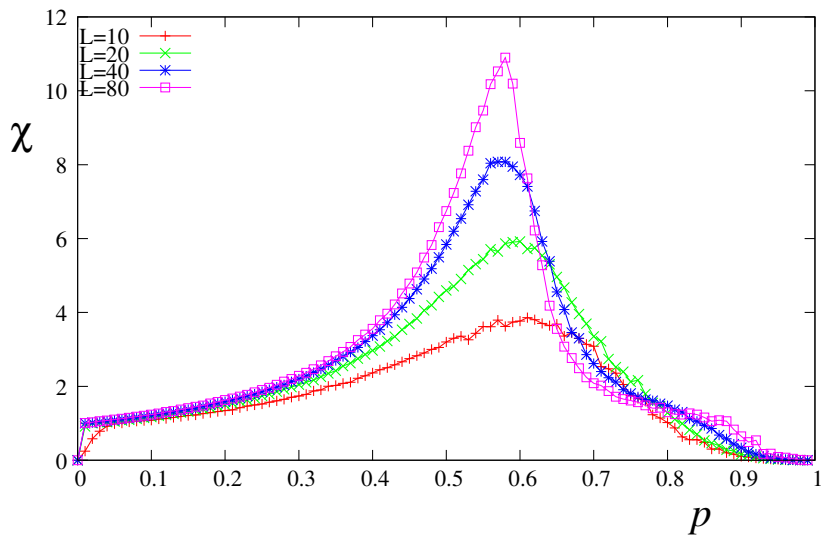
1. Is there an infinite cluster in infinite systems?
2. How many infinite clusters are there?
3. Cluster size distribution (n_s)
4. Mean cluster size (without the infinite one)? ($S = \sum_s s^2 n_s$)

Answers:

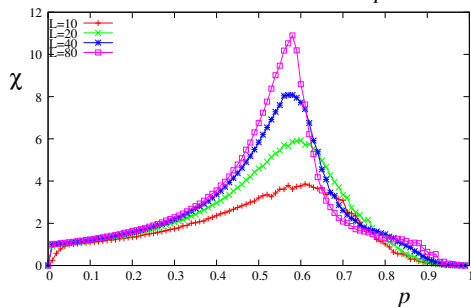
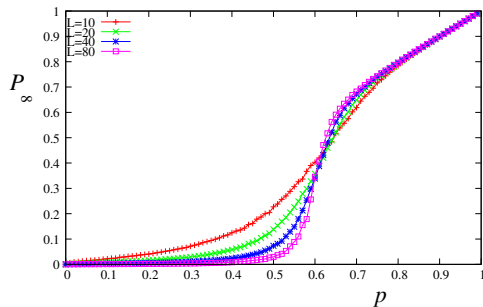
1. if $p > p_c$ then yes, otherwise no
2. Only 1!
3. $n_s \sim s^{-\tau}$
4. $S \sim |p - p_c|^{-\gamma}$

Like a second order phase transition in a geometric system!

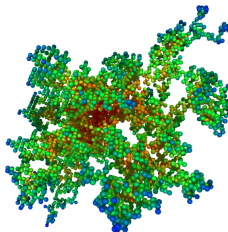
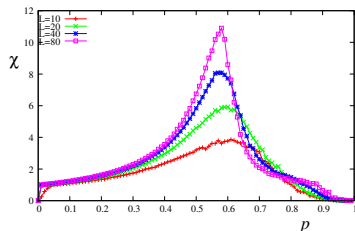
Percolation model



Percolation model



Percolating cluster



► Largest cluster

- fractal with fractal dimension of d_f

$$\text{► } S_\infty \sim \begin{cases} \xi_f^d \log(N/\xi_f^d) & p < p_c \\ N^{d_f/d} & p = p_c \\ NP_\infty(p) & p > p_c \end{cases}$$

- Largest not infinite cluster: size $\sim |p - p_c|^{-\nu}$

Percolation on networks (graphs)

- ▶ Network is defined by nodes and links
- ▶ Percolation gives us connected components
- ▶ Link removal percolation gives information about robustness, and structure

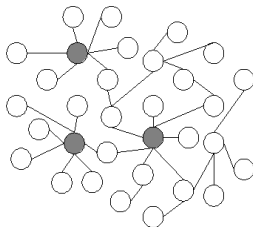
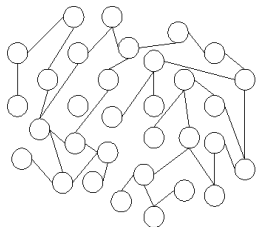


Percolation and attack on random networks

- ▶ Failure: equivalent to percolation: remove nodes at random
- ▶ Attack: remove most connected nodes

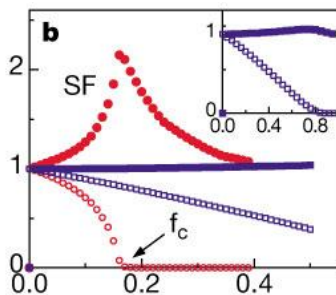
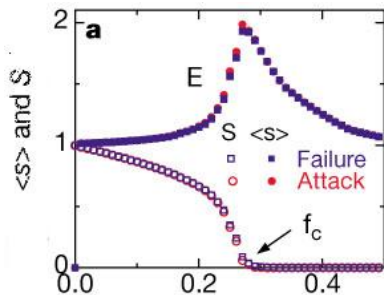


Error vs. attacks



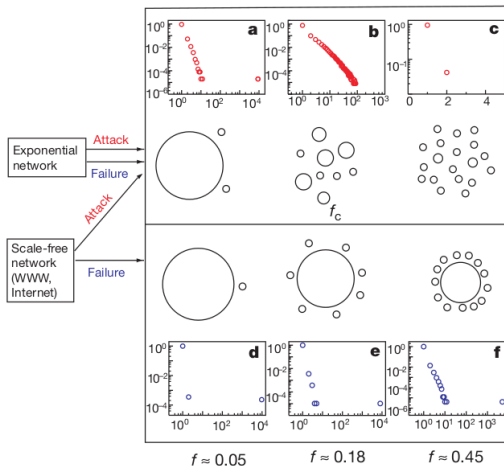
(a) Random network

(b) Scale-free network



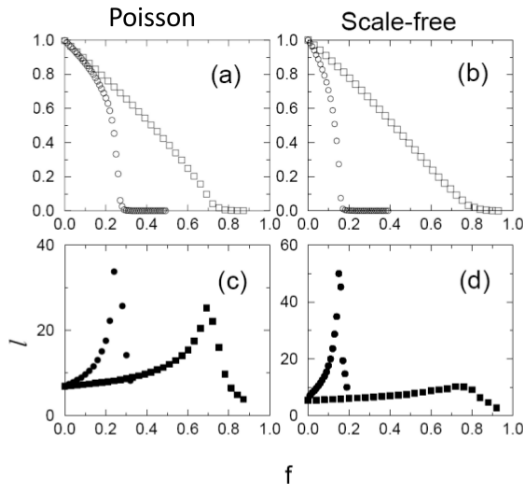
Percolation and attack on random networks

- ▶ Failure: equivalent to percolation: remove nodes at random
- ▶ Attack: remove most connected nodes



Robustness

- ▶ Link/node removal percolation
- ▶ Here: random, and largest first
- ▶ There is also weakest first



- squares: random failure
- circles: targeted attack

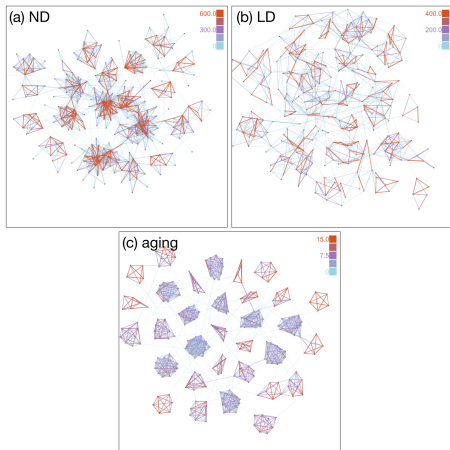
Failures: little effect on the integrity of the network if scale free.

Attacks: fast breakdown

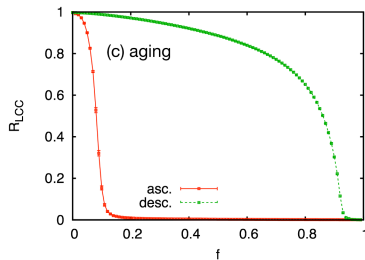
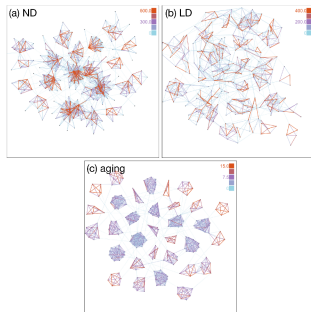
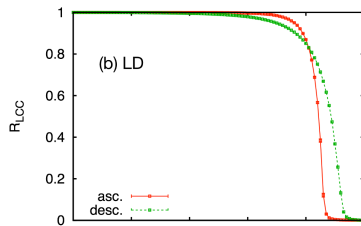
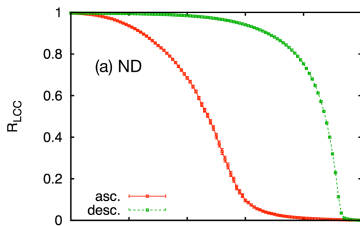
$\langle l \rangle$: average component size

Link removal percolation on networks

- ▶ Granovetter hypothesis: The strength of the weak ties
- ▶ Human communities have strong connections
- ▶ These communities are connected with weak ties
- ▶ Test the structures with Link removal percolation

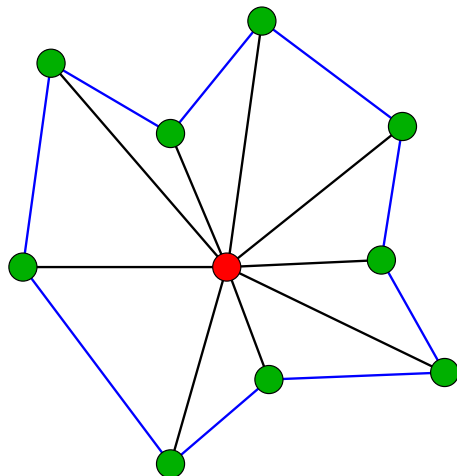


Link removal percolation on networks



Robustness

- ▶ Resistant both against random and targeted attacks.
- ▶ Must have hubs to resist random attacks
- ▶ Small degree nodes should be interconnected so they remain viable after removal of the hubs



Robustness against attacks

- ▶ Malicious attacks target central nodes, hubs
- ▶ Solution: central nodes should be connected
- ▶ Assortative mixing is preferred (high degree nodes are connected between each other)
- ▶ (Barabasi-Albert is thus a bad example)
- ▶ Robustness measure:

$$R = \frac{1}{N} \sum_{Q=1}^N s(Q)$$

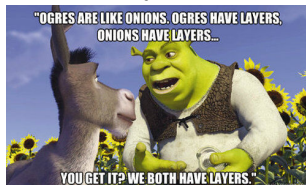
- ▶ $s(Q)$ fraction of nodes in the largest connected cluster after removing $Q = qN$ nodes
- ▶ Optimize for R

Onion structures

- ▶ Robustness measure:

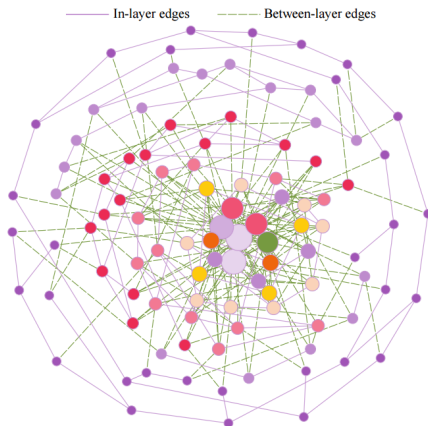
$$R = \frac{1}{N} \sum_{Q=1}^N s(Q)$$

- ▶ $s(Q)$ fraction of nodes in the largest connected cluster after removing $Q = qN$ nodes
- ▶ Optimize R by only rewiring and keeping degree distribution constant
- ▶ Onion structures are the most robust
 - ▶ Assortative
 - ▶ Layers with similar degree nodes
 - ▶ Inter-layer connections



Onion structures

- ▶ Assortative
- ▶ Layers with similar degree nodes
- ▶ Inter-layer connections

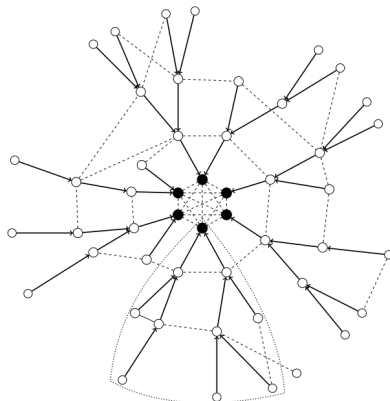


Château de Vincennes



Optimization: costs

- ▶ Internet Autonomous system topology
- ▶ Providers can connect to the top tier or be a customer
- ▶ They are responsible for directing the Internet traffic
- ▶ Simple protocols define the routing (mainly greedy)
- ▶ Many optimizes the structure



Flight route optimization

- ▶ Suppose weight of a link is defined as

$$w_{ij} = d_{ij}/t_{ij}$$

where d_{ij} is the distance, and t_{ij} is the traffic between two cities

- ▶ When more paths are possible the most economical is used:

$$C_{ij} = \min_{p \in \mathcal{P}} \sum_{l \in p} w_l$$

- ▶ Keep total traffic constant
- ▶ Function to be optimized is the average cost to pay to travel from any node to any other

$$\mathcal{L} = \frac{2}{N(N-1)} \sum_{i < j} C_{ij}$$

Flight route optimization

- ▶ Check a small circle:
- ▶ Let us assume $d_1 = d(A, B) = d(B, C) > d(A, C) = d'$
- ▶ Cost function (T is the average traffic between two cities):

$$\mathcal{L}_1 = \frac{2d + d'}{T}$$

- ▶ Cut connection (B, C) . The new cost function

$$\mathcal{L}_2 = \frac{d + d'}{2T} < \mathcal{L}_1$$

- ▶ The optimal path is a tree!

Tree model

- If it is known that the network is a tree task is easier:

$$\mathcal{L}_{\sqcup} = \sum_{e \in T} b_e \frac{d_e}{t_e}$$

where b_e is the link betweenness centrality

- The optimal traffic

$$t_e = \frac{T \sqrt{b_e d_e}}{\sum_e \sqrt{b_e d_e}}$$

- The optimal traffic tree can then be obtained by minimizing

$$\mathcal{L} = \sum_{e \in T} \sqrt{b_e d_e}$$

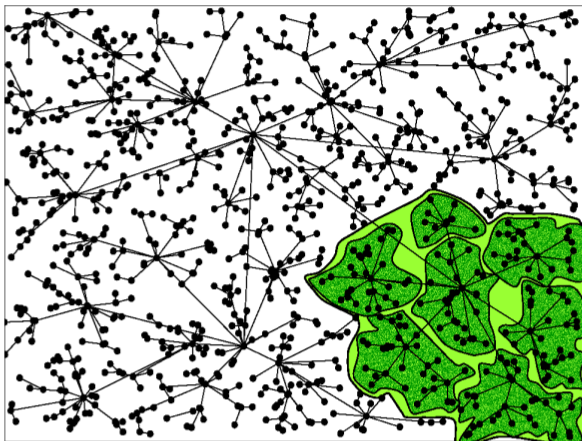
- More generally

$$\mathcal{L} = \sum_{e \in T} b_e^\mu d_e^\nu$$

where μ and ν control the relative importance of distance against topology as measured by centrality

Optimal traffic on networks

- ▶ Exponential degree distribution
- ▶ Power law betweenness distribution
- ▶ Hierarchical organizations
- ▶ $\mu = \nu = 0.5$



Optimal traffic on networks

