

# Computer Simulations in Physics

## Course for MSc physics students

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# Information

## ▶ Coordinates:

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▶ Consultation:

▶ F III building, first floor 6 (after the first stairs to the right, at the end of the corridor), Department of Theoretical Physics

▶ Teams or Meet

▶ Upon demand (Email)

## ▶ Webpage:

[http://physics.bme.hu/BMETE15MF74\\_kov?language=en](http://physics.bme.hu/BMETE15MF74_kov?language=en)

▶ Homework: <http://edu.ttk.bme.hu/>

## Required knowledge

- ▶ Knowledge of basic statistical physics
- ▶ Knowledge of basic quantum mechanics
- ▶ C, C#, C++ or python language
- ▶ If you use C# please submit only the code part!

# Potential skills earned

- ▶ Learn simulation methods
- ▶ Implementing models
- ▶ Optimization
- ▶ Use chatGPT

# Programming environment

- ▶ Basic
  - ▶ Editor (not notepad!)
  - ▶ Compiler (gcc recommended)
- ▶ Advanced
  - ▶ Developer environment
  - ▶ Integrated developer environment (also compiles)
- ▶ Super advanced
  - ▶ ChatGPT, use it and show how it can be used what are the limitations etc.

# Integrated developer environment

- ▶ Visual studio (old version can be downloaded from <http://software.eik.bme.hu> only from bme.hu domain)
- ▶ Eclipse
- ▶ Netbeans
- ▶ CodeLite
- ▶ pyCharm
- ▶ Anaconda
- ▶ Google colab
- ▶ etc.

# Install compiler

## ▶ Linux

- ▶ Install development package, usually not installed when desktop installation was selected (libgcc-version-dev, plus any -dev packages you want)

## ▶ Windows

- ▶ Visual studio
- ▶ cygwin+gcc <http://preshing.com/20141108/how-to-install-the-latest-gcc-on-windows/>
- ▶ Eclipse+gcc (eclipse does not come with a C compiler)  
[http://www.dcs.vein.hu/bertok/oktatas/cpp\\_by\\_eclipse/eclipse\\_for\\_cpp\\_on\\_windows.html](http://www.dcs.vein.hu/bertok/oktatas/cpp_by_eclipse/eclipse_for_cpp_on_windows.html)
- ▶ Anaconda  
<https://www.anaconda.com/products/individual>
- ▶ Linux in Virtualbox

# Random numbers

- ▶ Gnu Scientific library
  - ▶ variable: `gsl_rng *r;`
  - ▶ reading environment variables `GSL_RNG_TYPE` and `GSL_RNG_SEED`: `gsl_rng_env_setup`
  - ▶ `gsl_rng_default=gsl_rng_mt19937` Mersenne twister algorithm period:  $2^{19937} - 1$
  - ▶ Set seed: `gsl_rng_set(r,seed);`
  - ▶ Integer random numbers between `gsl_rng_max(r)` and `gsl_rng_min(r)`: `unsigned long gsl_rng_get(r);`
  - ▶ *double* random numbers from 0 to 1 (0 included, 1 excluded) `gsl_rng_uniform(r);`



# Requirements

- ▶ **Minimum requirements**
  - ▶ 5 homeworks submitted and accepted ( $> 50\%$ )
  - ▶ Presented and accepted project
- ▶ **Exam: mark**
  - ▶ **500 point:** 5 homeworks (deadlines are **soft!**)
  - ▶ **400 points:** Small (30 min) test (individual) (must pass!) **or** 4 extra homeworks
  - ▶ **900 points:** From projects (pairs/groups) presented at the end of the course.
  - ▶ **20-50 points/piece** extra: practice exercises can earn you 20 extra points each. Deadlines are **hard!**
  - ▶ The marks will be calculated using the 1800 point threshold
  - ▶ Upon request please, be ready for a code checkup to verify ownership. During this check you will be shown parts of your code and you are supposed to explain wht it was meant to do.
  - ▶ Turn it in language: English, Hungarian, German, French

# Requirements

## ▶ Homeworks

- ▶ Individual work. You can use the Internet or chatGPT, but then explain what part of the code came where and what modifications were needed. A good comparison of a chatGPT code with an improved own one earns you extra points
- ▶ Documented working codes (no extra libraries except for `gs1`, `numpy`, `scipy`)
- ▶ Always include an image from the results or some documentation.
- ▶ Please combine all files in a zip/rar/tgz/etc. archive and upload a single file
- ▶ Python notebooks are also accepted, please do not clear the results from the notebook!
- ▶ Using fancy visualization techniques does not improve the mark which is given for the algorithm, the efficiency of the code and the solution of the problem
- ▶ Please, keep in mind that I do not have time to compile and run your code. Make it human readable!

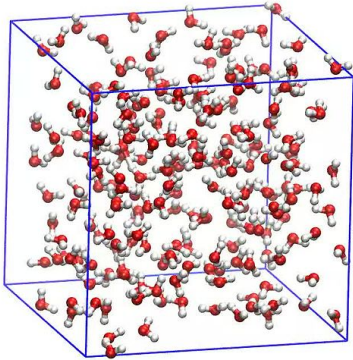
# Literature

- ▶ There used to be a list of outdated books here.
- ▶ I will try to give you material on subject basis: For this one:
- ▶ `http://www.lce.hut.fi/teaching/S-114.1100/lect_8.pdf`
- ▶ `https://arxiv.org/pdf/1005.4117.pdf`
- ▶ `https://www.ks.uiuc.edu/Services/Class/PHYS498NSM/`

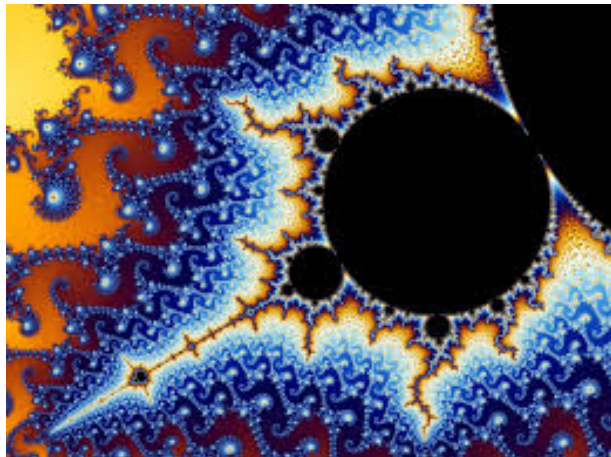
# Subjects

1. Random numbers
2. Molecular dynamics
3. Other particle based methods
4. Percolation, Fractals
5. Ising, Schelling
6. Schrödinger equation
7. Quantum scattering
8. Optimization (annealing, genetic)
9. Complex networks
10. Clustering, community detection
11. Algorithmically defined models
12. Neural networks
13. Game models
14. Presentation

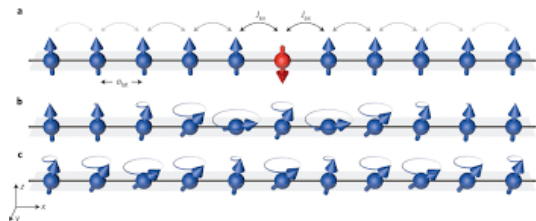
# Discrete element methods



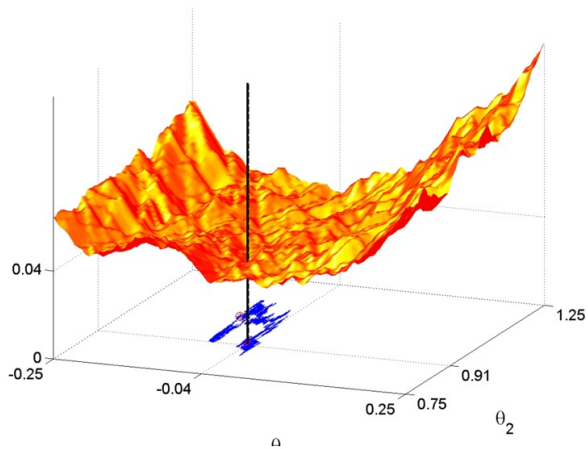
# Percolation, Fractals



# Ising, Heisenberg model

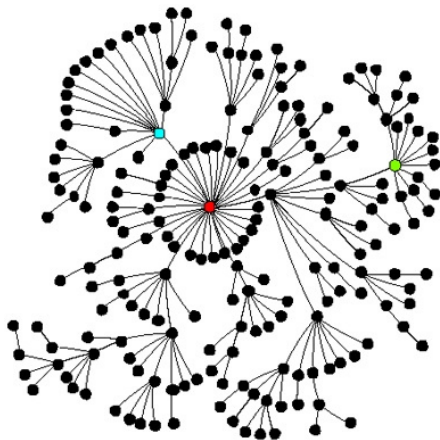


# Optimization

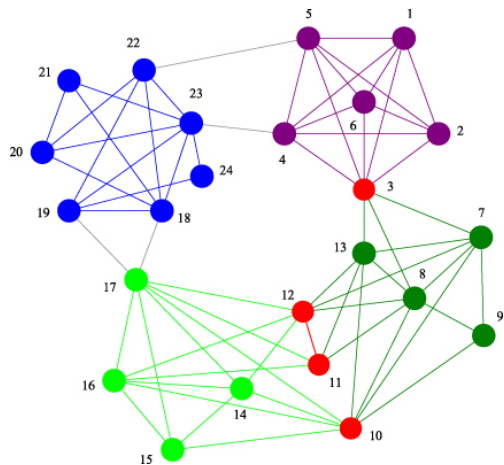




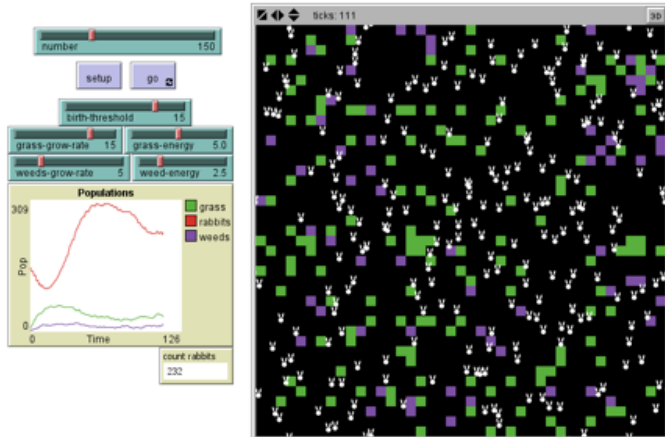
## Complex networks



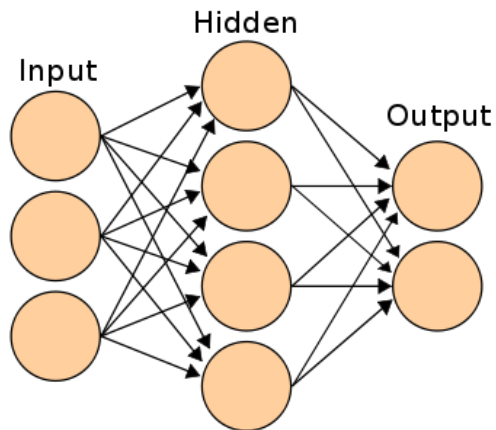
# Clustering



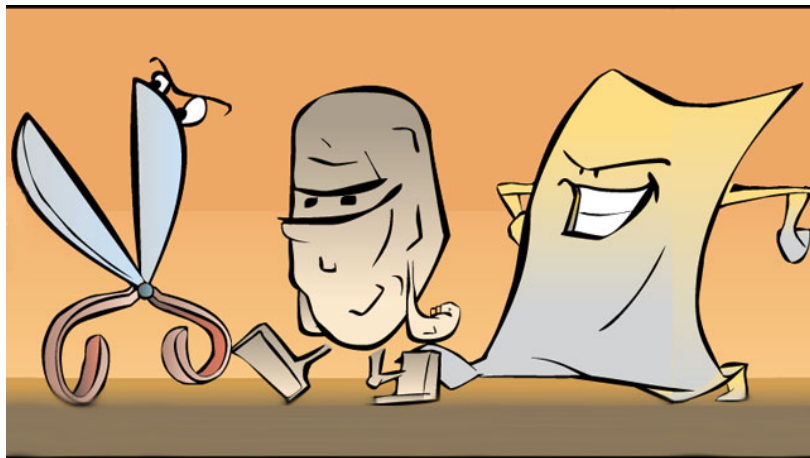
# Algorithmically defined models



# Neural networks



# Game models



# Simulations

## *Experiments*

Principle of measurement

Apparatus

Calibration

Sample

Measurement

## *Simulations*

Algorithm

Program + Hardware

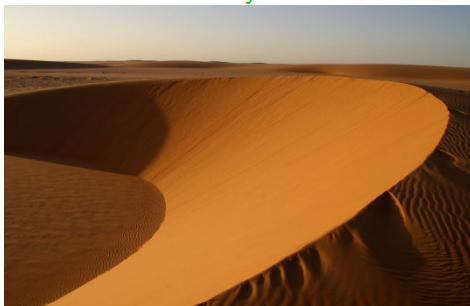
Calibration + Debugging

Sample

Run

Data collection

Analysis



# Programming languages

## Simulations codes

- ▶ System size must be large
  - ▶ Phase transition  $\xi \rightarrow \infty$
  - ▶ Real systems  $N \sim 10^{23}$  (memory  $< 10^{12}$ )
- ▶ Simulation time should be long
  - ▶ Relaxation time
  - ▶ Interesting phenomena take long
  - ▶ Separation of time scales

Must be efficient!

It is not bad if program is readable and extensible...

## Sample preparation

- ▶ Sometimes it is also a simulation

## Data analysis

- ▶ Anything may happen

# Programming languages

## Problem to solve:

- ▶ Fill an array with product of two random numbers
- ▶ Calculate the average of them

### python

```
#!/usr/bin/python
import random

N = 20000000
s = []

for a in range(N):
    s.append(random.random() * random.random())

av = 0
for a in range(N):
    av += s[a]

print av / N
```

### C

```
#include <stdlib.h>
#include <stdio.h>
#include <math.h>

int main(np,para)
char *para[];
{
    int N = 200000000;
    int a;
    double *s,av;

    s = (double *) malloc(sizeof(double) * N);
    av = 0.0;
    for (a=0;a<N;a++) {
        s[a] = (double)rand() / RAND_MAX * rand() / RAND_MAX;
    }
    for (a=0;a<N;a++) {
        av += s[a];
    }
    printf("%lg\n",av/N);
}
```



# Programming languages

```
#!/usr/bin/python
```

```
import random
```

```
N = 20000000
```

```
s = []
```

```
for a in range(N):
```

```
    s.append(random.random() * random.random())
```

```
av = 0
```

```
for a in range(N):
```

```
    av += s[a]
```

```
print av / N
```

```
#!/usr/bin/python
```

```
import numpy
```

```
N = 200000000
```

```
s = numpy.random.random(N)
```

```
s *= numpy.random.random(N)
```

```
print s.mean()
```

6.9s	3.46s
4.51s	3.29s

```
#include <stdlib.h>
#include <stdio.h>
#include <math.h>
```

```
int main(np,para)
char *para[];
```

```
{
    int N = 200000000;
    int a;
    double *s,av;
```

```
    s = (double *) malloc(sizeof(double) * N);
    av = 0.0;
```

```
    for (a=0;a<N;a++) {
        s[a] = (double)rand() / RAND_MAX * rand() / RAND_MAX;
    }
```

```
    for (a=0;a<N;a++) {
        av += s[a];
    }
```

```
    printf("%lg\n",av/N);
```

```
}
#include <stdlib.h>
#include <stdio.h>
#include <math.h>
```

```
int main(np,para)
char *para[];
```

```
{
    int N = 200000000;
    int a;
    double *s,av;
```

```
    s = (double *) malloc(sizeof(double) * N);
    av = 0.0;
```

```
    for (a=0;a<N;a++) {
        s[a] = (double)rand() / RAND_MAX * rand() / RAND_MAX;
        av += s[a];
    }
```

```
    printf("%lg\n",av/N);
```

```
}
```

# Optimization

- ▶ **Programming language**
  - ▶ In example C is 1.3-2 times faster than python
  - ▶ Matlab can be very efficient, but it is proprietary
  - ▶ Matlab, Maple, Mathematica are expensive
  - ▶ All clusters have C, and C++, and python
- ▶ **Optimization**
  - ▶ Parallelization
  - ▶ Indexing, careful usage of pointers
  - ▶ Reformulate operations
  - ▶ Does not always worth the pain
  - ▶ gprof

Flat profile:

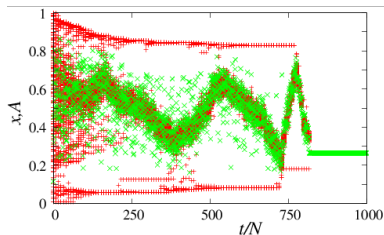
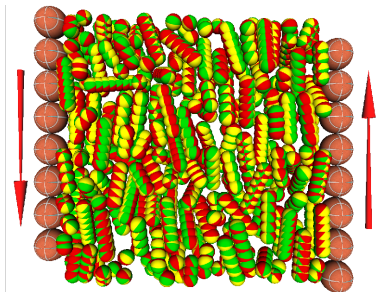
Each sample counts as 0.01 seconds.

% time	cumulative seconds	self seconds	calls	self ms/call	total ms/call	name
37.66	56.83	56.83	3248064862	0.00	0.00	is_in_community
25.99	96.05	39.22	1000000	0.04	0.04	e_erode
11.55	113.47	17.43	21355853	0.00	0.00	weighted_random_link
6.33	123.03	9.55	11078805	0.00	0.00	weighted_random_link_ban_list
3.02	127.58	4.55	8406648	0.00	0.01	e_info
2.77	131.75	4.18				main
2.26	135.16	3.40	197988614	0.00	0.00	ct_weight
2.10	138.33	3.17	4	792.50	792.50	clear_data
1.85	141.12	2.79	12949626	0.00	0.00	e_single
1.73	143.74	2.62	164260875	0.00	0.00	ranksz
1.60	146.16	2.42	12774907	0.00	0.00	strengthen
0.97	147.62	1.46	19359356	0.00	0.01	communicate
0.88	148.94	1.33	248428917	0.00	0.00	is_internet
0.32	149.43	0.48	15380	0.03	0.03	random_agent_with_group_sex
0.31	149.90	0.47	2042439	0.00	0.00	e_share
0.24	150.25	0.36				seed3

---

# Simulations

- ▶ Do what nature does
  - ▶ Molecular dynamics
  - ▶ Hydrodynamics
- ▶ Make use of statistical physics
  - ▶ Monte-Carlo dynamics
  - ▶ Simulate simplified models
  - ▶ Much smaller codes!



# Random numbers

- ▶ Why?
  - ▶ Ensemble average:

$$\langle A \rangle = \sum_i A_i P_i^{\text{eq}}$$

Random initial configurations

- ▶ Model: e.g. Monte-Carlo
  - ▶ Fluctuations
  - ▶ Sample
- ▶ How?





# Random number generators

- ▶ True (Physical phenomena):
  - ▶ Shot noise (circuit)
  - ▶ Nuclear decay
  - ▶ Amplification of noise
    - ▶ Atmospheric noise (random.org)
    - ▶ Thermal noise of resistor
    - ▶ Reverse biased transistor
    - ▶ Lava lamps
  - ▶ Limited speed
  - ▶ Needed for cryptography
- ▶ Pseudo (algorithm):
  - ▶ Deterministic
    - ▶ Good for debugging!
  - ▶ Fast
  - ▶ Can be made reliable



# Language provided random numbers

It is good to know what the computer does!

- ▶ Algorithm
  - ▶ Performance
  - ▶ Precision
  - ▶ Limit cycle
  - ▶ Historically(?) a catastrophe

```
int getRandomNumber()  
{  
    return 4; // chosen by fair dice roll.  
             // guaranteed to be random.  
}
```



# Language provided random numbers

It is good to know what the computer does!

```
int getRandomNumber()  
{  
    return 4; // chosen by fair dice roll.  
             // guaranteed to be random.  
}
```

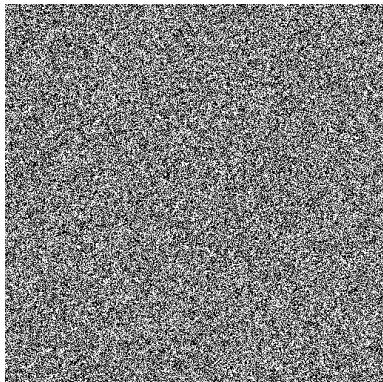
**DILBERT** By SCOTT ADAMS



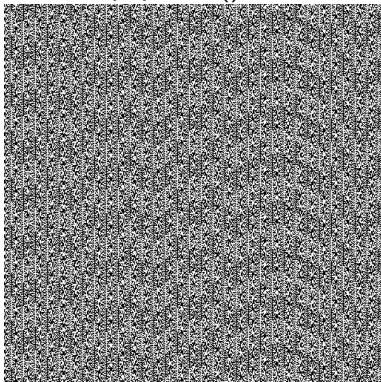
# Language provided random numbers

It is good to know what the computer does!

Random



php rand() on Windows



# Language provided random numbers

It is good to know what the computer does!

- ▶ Algorithm
  - ▶ Performance
  - ▶ Precision
  - ▶ Limit cycle
  - ▶ Historically a catastrophe
- ▶ Seed
  - ▶ From true random source
  - ▶ Time
  - ▶ **Manual**
    - ▶ Allows debugging
    - ▶ Ensures difference

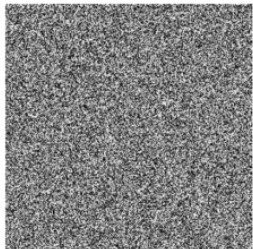
First only uniform random numbers

# Seed

- ▶ From true random source
- ▶ Time
- ▶ **Manual**

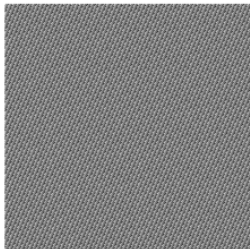
Random number generator of Python with different seeds:

System.Random  
numbers 0...n of seed 0



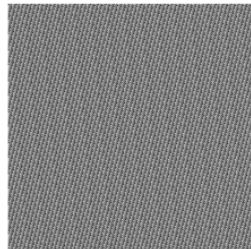
Sequence of 65536 random values.

System.Random  
0th number of seed 0...n



Sequence of 65536 random values.

Linear function  $i * 19969 / 207$   
numbers 0...n

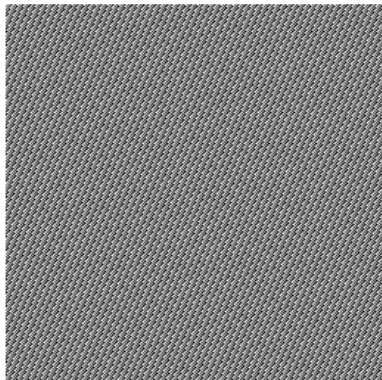


Sequence of 65536 random values.

## Seed

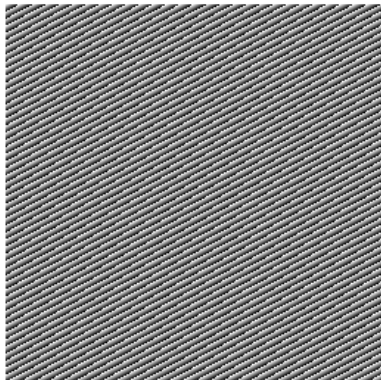
- ▶ Ensemble average: Include in the code if possible instead of restarting it with different seeds!

**System.Random**  
0th number of seed 0...n



Sequence of 65536 random values.

**System.Random**  
100th number of seed 0...n



Sequence of 65536 random values.

# Multiplicative congruential algorithm

- ▶ Let  $r_j$  be an integer number, the next is generated by

$$r_{j+1} = (ar_j + c) \bmod(m),$$

- ▶ Sometimes only  $k$  bits are used
- ▶ Values between 0 and  $m - 1$  or  $2^k - 1$
- ▶ Three parameters  $(a, c, m)$ .
- ▶ If  $m = 2^X$  is fast. Use AND (&) instead of modulo (%).
- ▶ Good:
  - ▶ Historical choice:  
 $a = 7^5 = 16807$ ,  $m = 2^{31} - 1 = 2147483647$ ,  $c = 0$
  - ▶ gcc built-in ( $k = 31$ ):  
 $a = 1103515245$ ,  $m = 2^{31} = 2147483648$ ,  $c = 12345$
- ▶ Bad:
  - ▶ RANDU:  $a = 65539$ ,  $m = 2^{31} = 2147483648$ ,  $c = 0$

# Tausworth, Kirkpatrick-Stoll generator

- ▶ Fill an array of 256 integers with random numbers

$$J[k] = J[(k - 250) \& 255] \wedge J[(k - 103) \& 255]$$

- ▶ Return  $J[k]$ , increase  $k$  by one
- ▶ Can be 64 bit number
- ▶ Extremely fast, but short cycles for certain seeds

XOR function

$\wedge$	1	0
1	0	1
0	1	0

# Tausworth, Kirkpatrick-Stoll generator corrected by Zipf

## The one the lecturer uses

- ▶ Fill an array of 256 integers with random numbers

$$J[k] = J[(k - 250) \& 255] \wedge J[(k - 103) \& 255]$$

Increase  $k$  by one

$$J[k] = J[(k - 30) \& 255] \wedge J[(k - 127) \& 255]$$

- ▶ Return  $J[k]$ , increase  $k$  by one
- ▶ Extremely fast, reliable also on bit level
- ▶ General transformation  $x \in [0 : 1[$

$$x = r / (RAND\_MAX + 1)$$



## Floating point random numbers

- ▶ General transformation  $x \in [0 : 1[$

$$x = r / (RAND\_MAX + 1)$$

- ▶ It is important to know whether limits are included or not
- ▶ General feature: 0 included 1 not
- ▶ Generate integer number from 1,2,3. use  $i = r \% 3$  (modulo)  
result: 1 will be  $1 + 10^{-9}$  more probable than 2 or 3.
- ▶ General practice use division instead of percentage, higher bits are more reliable for LCG

# Tests

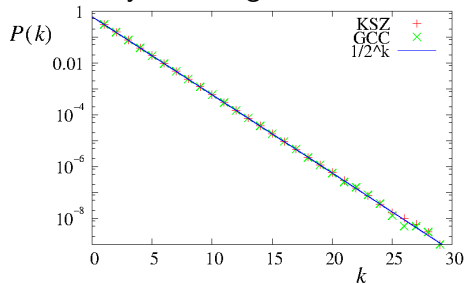
- ▶ General: e.g. TESTU01
- ▶ Diehard tests:
  - ▶ Birthday spacings (spacing is exponential)
  - ▶ **Monkey tests (random typewriter problem)**
  - ▶ Parking lot test
  - ▶ Moments:  $m = \int_0^1 \frac{1}{n+1}$
  - ▶ Correlation  
$$C_{q,q'}(t) = \int_0^1 \int_0^1 x^q x'^{q'} P[x, x'(t)] dx dx' = \frac{1}{(q+1)(q'+1)}$$
  - ▶ Fourier-spectra
  - ▶ **Fill of  $d$  dimensional lattice**
  - ▶ **Random walks**

**Red ones are not always fulfilled!**

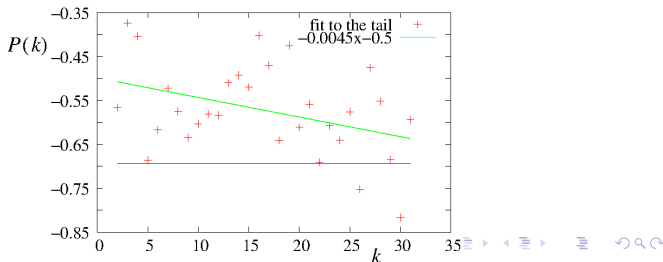
- ▶ Certain Multiplicative congruential generators are bad on bit series distribution, not completely position independent.

## Bit series distribution

Probability of having  $k$  times the same bit

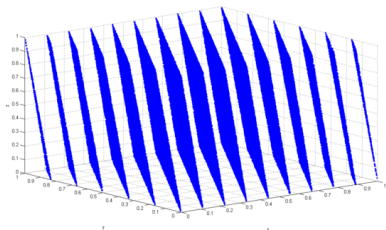


Fit to the tail for different bit positions (gcc)



## Fill of $d$ dimensional lattice

- ▶ Generate  $d$  random numbers  $c_i \in [0, L]$
- ▶ Set  $x[c_1, c_2, \dots, c_d] = 1$
- ▶ The Marsaglia effect is that for all congruential multiplicative generators there will be unavailable points (on hyperplanes) if  $d$  is large enough.
- ▶ For RANDU  $d = 3$



## Solution for Marsaglia effect

- ▶ Instead of  $d$  random numbers only 1 ( $x$ )
- ▶ Divide it into  $d$  parts:  $k = \text{int}(\log_d(\text{RANDMAX}))$   
 $c_1 = x \% k, x /= k$   
 $c_2 = x \% k, x /= k$   
...
- ▶ Better to have  $L = 2^k$ . Which is much faster because of AND and SHIFT operations

General advice: Save time by generating less random numbers

## Random numbers with different distributions

- ▶ Let us have a good random number  $r \in [0, 1]$ .
- ▶ The probability density function is  $P(x)$
- ▶ The cumulative distribution is

$$D(x) = \int_{-\infty}^x P(x') dx'$$

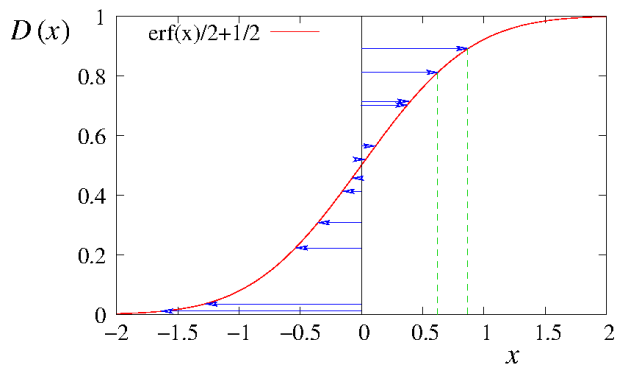
- ▶ Obviously:

$$P(x) = D'(x)$$

- ▶ The numbers  $D^{-1}(x)$  will be distributed according to  $P(x)$
- ▶  $D^{-1}(x)$  is the inverse function of  $D(x)$  not always easy to get!

# Random numbers with different distributions

## Graphical representation



# Random numbers with different distributions

A soluable example



$$P(x) = \frac{1}{\pi} \frac{1}{1+x^2}$$



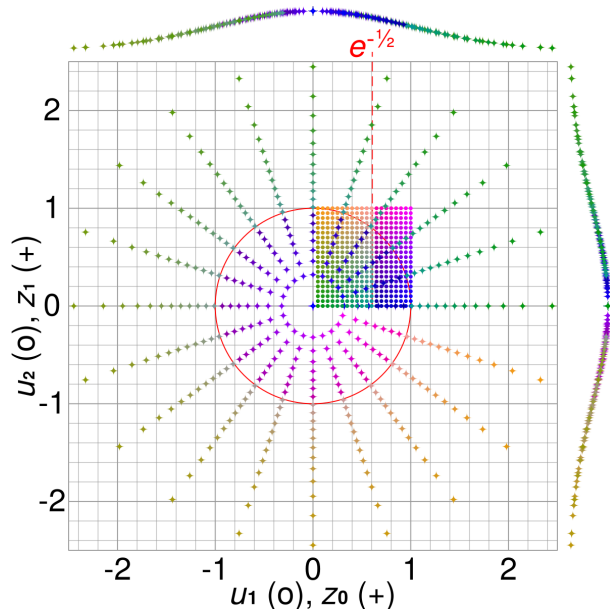
$$D(x) = \frac{1}{\pi} \int_{-\infty}^x \frac{1+x'^2}{d} x' = \frac{1}{2} + \frac{1}{\pi} \arctan(x)$$



$$x = D^{-1}(y) = \tan [\pi(y - 1/2)]$$



# Box-Müller method



## Box-Müller method

- ▶  $r_1, r_2$  uniformly distributed between 0 and 1
- ▶ if  $r_1^2 + r_2^2 \geq 1$  or  $r_1 = r_2 = 0$  discard  $r_1$  and  $r_2$ . So the generated point is inside the unit circle and is not at the origin
- ▶  $R^2 = r_1^2 + r_2^2$  is uniformly distributed between 0 and 1.
- ▶  $U_1 \equiv R^2, U_2 \equiv \arctan(r_1/r_2)/(2\pi)$
- ▶ Two independent normally distributed random numbers:

$$x_1 = \sqrt{-2 \log U_1} \cos(2\pi U_2) = r_1 \sqrt{\frac{-2 \log(R^2)}{R^2}}$$

$$x_2 = \sqrt{-2 \log U_1} \sin(2\pi U_2) = r_2 \sqrt{\frac{-2 \log(R^2)}{R^2}}$$

## Power law distributed random numbers

Let  $P(y)$  have uniform distribution in  $[0, 1]$ . We generate  $P(x)$  such as

$$P(x) = Cx^n$$

for  $x \in [x_0, x_1]$ .

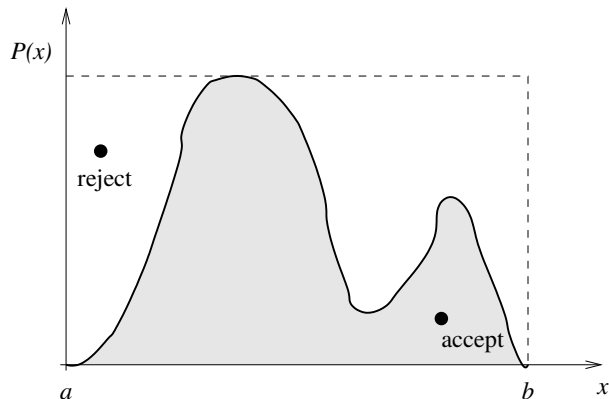
$$D(x) = \int_{x_0}^x P(x') dx' = \frac{C}{n+1} (x^{n+1} - x_0^{n+1})$$

The inverse function is simple:

$$x = [(x_1^{n+1} - x_0^{n+1})y + x_0^{n+1}]^{1/(n+1)}$$

# Monte Carlo

- ▶ Identify base:  $[a, b]$
- ▶ Identify minimum/maximum:  $P_{\max} = \max_{x \in [a, b]} P(x)$ , idem...
- ▶ Generate a point  $(x, y)$  in the rectangle  $(a, P_{\min}), (b, P_{\max})$
- ▶ If  $y < P(x)$  the return  $x$  otherwise generate new point



## Error

- ▶ Ensemble average
- ▶ Example: estimate  $\pi$
- ▶ Drop a needle of length  $l \leq t$
- ▶ May or may not cross a line

$$p_{\text{cross}} = \frac{2l}{t\pi}$$

- ▶ Lazzarini in 1901 using  $N = 3408$  tries got:

$$\pi \simeq 355/113 = 3.14159292 = \pi + \mathcal{O}(10^{-7})$$

- ▶ Impressive  $10^{-7}$  error, but

$$\frac{1}{\sqrt{N}} \simeq 0.0017$$

